# Marginal Waiting Cost in Optimization Based Flow Control

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### Flow control (congestion control)



#### Example:

TCP flow control & random early detection (RED)

- *p*: Marking in an ACK packet
- *x<sub>i</sub>* corresponds to congestion window size.

ABR flow control

*p*: Congestion indication (CI) and/or explicit rate (ER) in an RM cell r corresponde to ellowed cell rate (ACR)

*x<sub>i</sub>* corresponds to allowed cell rate (ACR)

### Flow control as an optimization problem

To construct a flow control mechanism

- --> To determine the following according to a policy
  - How to generate p
  - How to control  $x_i$  according to p



- Policy

- --> To minimize the total cost, which is congestion cost minus user utilities
- Mechanism to generate p and to control  $x_i$ 
  - --> Solution for the optimization problem corresponding to the policy

## Related works

F. P. Kelly et al., Rate control for communication networks: shadow price, proportional fairness and stability, Journal of the Operational Research Society, 49 (1998).

- S. J. Golestani et al., A class of end-to-end congestion control algorithms, Sixth International Conference on Network Protocols (1998).

- S. H. Low et al., Optimization flow control, I: basic algorithm and convergence, IEEE/ACM Transactions on Networking (1999).

Kelly et al. (1998)

- Policy (single bottleneck case)

 $\max_{\mathbf{x}} J(\mathbf{x}) = \sum_{i} U_{i}(x_{i}) - R(\sum_{i} x_{i})$ 

 $U_i(\cdot)$ : Utility function of user *i*, strictly increasing, concave, differentiable  $R(\cdot)$ : Congestion cost of the bottleneck link, differentiable p(y) = d/dy R(y): shadow price, positive, strictly increasing

- Mechanism to control  $x_i$  (willingness to pay type control)

$$d/dx \ x_i(t) = \kappa_i \left( w_i(t) - x_i(t) p\left(\sum_j x_j(t)\right) \right)$$

 $w_i(t) = x_i(t)U'(x_i(t))$ 

*κi*: control parameter

- To consider accumulated cost, for example,

Minimize  $\int \left( \sum_{i} U_{i}(x_{i}(t)) - R\left( \sum_{i} x_{i}(t) \right) \right) dt$ 

- To consider stochastic models, which are queueing models

#### Single bottleneck model: single node and multiple users



- Poisson arrival processes:
  - intensity vector  $\lambda(t) = (\lambda_i(t))$  <-- control subject
- Exponential distributed services: service rate  $\mu$  (the same for all the users)
- *K*: Buffer size, including service facility
- L(t): the number of packets in the system at time t
- $A_i(t)$ : the number of user *i*'s packets arriving in [0, t]

## User utility and congestion cost

User utility is represented as a function of throughput.

 $U_i(1(L(t) < K) \lambda_i(t))$ : **Instantaneous utility** of user *i* at time *t* 

 $U_i(\cdot)$ : a function, non-negative, increasing, differentiable, strictly concave

$$\mathbf{E}\left[\frac{1}{t}\int_{0}^{t}\mathbf{1}(L(t-) < K)dA_{i}(t)\right] = \mathbf{E}\left[\frac{1}{t}\int_{0}^{t}\mathbf{1}(L(t) < K)\lambda_{i}(t)dt\right]$$

 $1(\cdot)$ : indicator function

--> From this formula,  $1(L(t) < K) \lambda_i(t)$  can be considered as **the instantaneous throughput** of user *i* 

Congestion cost approximately represents delay and loss.

 $R(L(t)) \lambda_i(t)$ : Instantaneous congestion cost for user *i* a time *t* 

 $R(\cdot)$ : a function, non-negative, increasing (the same for all the users)

**Note**:  $R(K) \lambda_i(t)$  represents the instantaneous cost of loss.

# Formulation 1: Expected discounted cost

#### To minimize the expectation of the discounted total user cost

 $\mathbf{u}(t) = (ui(t))$ : Control parameters, i.e.,  $\lambda i(t) < -- ui(t)$ Bi = [bi,0, bi,1]: Range of ui(t),  $\mathbf{B} = B_1 \times ... \times B_k$  $U_{i,\max} = Ui(bi,1) b_{i,1}$ 

 $\alpha > 0$ : Exponential discount factor

Instantaneous cost for user *i* 

$$\begin{split} C_i(L(t),\lambda_i(t)) &\equiv \left\{ U_{i,\max} - U_i \Big( \mathbb{1}(L(t) < K)\lambda_i(t) \Big) \right\} + R(L(t))\lambda_i(t) \\ C(L(t),\lambda(t)) &\equiv \sum_{i=1}^k C_i(L(t),\lambda_i(t)) \end{split}$$

- Objective function  $J_{\alpha}(\mathbf{u}, l_{0}) \equiv \lim_{T \to \infty} \mathbb{E}_{\mathbf{u}} \left[ \int_{0}^{T} e^{-\alpha t} C(L(t), \mathbf{u}(t)) dt \mid L(0) = l_{0} \right]$ 

Policy (criterion) Minimize  $J_{\alpha}(\mathbf{u}, l_0)$ 

### Result 1: Expected discounted cost

Problem 1: P1  $J_{\alpha}^{*}(l_{0}) \equiv \inf_{\mathbf{u}} J_{\alpha}(\mathbf{u}, l_{0})$ 

#### Result

Assume that there exist the function  $V\alpha(\cdot)$  that satisfies

(1) 
$$V_{\alpha}(l) \equiv \inf_{\mathbf{u}} J_{\alpha}(\mathbf{u}, l),$$

and the function  $v * \alpha(\cdot)$  that satisfies

(2) 
$$\mathbf{v}_{\alpha}^{*}(l) = \underset{\mathbf{v}\in\mathbf{B}}{\operatorname{argmin}} \left\{ \sum_{i=1}^{k} v_{i} 1(l < K) \left[ V_{\alpha}(l+1) - V_{\alpha}(l) \right] + C(l,\mathbf{v}) \right\}.$$

Then, the optimal control  $\mathbf{u}^{*\alpha}$  is given by

(3) 
$$\mathbf{u}_{\alpha}^{*}(t) = \mathbf{v}_{\alpha}^{*}(L(t)).$$

This result is directly derived from Theorem VTT-T1 in Point Processes and Queues by Bremaud.

### Formulation 2: Expected average cost

To minimize the expectation of the average total user cost





### Result 2: Expected average cost

Problem 2: P2  
$$J^*(l_0) \equiv \inf_u J(u, l_0)$$

#### Result

Assume that there exist the function  $V\alpha(\cdot)$  that satisfies

(1) 
$$V_{\alpha}(l) = \inf_{u} J_{\alpha}(\mathbf{u}, l),$$

and the function  $G(\cdot)$  that is a certain limit of  $V\alpha(.)$ , i.e,

(4) 
$$G(l) = \lim_{n \to \infty} \left( V_{\alpha_n}(l+1) - V_{\alpha_n}(l) \right)$$
, where  $\alpha_n \to 0$  as  $n \to \infty$ .

Furthermore, assume that there exists the function  $v^*(\cdot)$  that satisfies

(5) 
$$\mathbf{v}^*(l) = \underset{\mathbf{v} \in \mathbf{B}}{\operatorname{argmin}} \left\{ \sum_{i=1}^k v_i \mathbf{1}(l < K) G(l) + C(l, \mathbf{v}) \right\}.$$

Then, the optimal control **u**\* is given by

(6) 
$$\mathbf{u}^*(t) = \mathbf{v}^*(L(t))$$

### Discussion

Functions  $V_{\alpha}(\cdot)$  and  $G(\cdot)$   $V_{\alpha}(\cdot)$  is the expected total cost for the future.  $V_{\alpha}(l) \equiv \inf_{\mathbf{u}} J_{\alpha}(\mathbf{u}, l)$   $G(\cdot)$  can be considered as the expected cost of one packet for the future.  $G(l) \equiv \lim_{n \to \infty} \left( V_{\alpha_n}(l+1) - V_{\alpha_n}(l) \right)$ 

Flow control (the case of expected average cost)

From equation (5),  $v^*(\cdot)$  satisfies

7) 
$$U'_i(v^*_i(l)) \equiv G(l) + R(l), \ l < K.$$

Therefore, signal  $p(\cdot)$  can be define as

 $p(l) \equiv G(l) + R(l),$ 

and the packet arrival intensities are given as the solution of

P3) 
$$\inf_{v_i} \{ v_i p(L(t)) - U_i(v_i) \}.$$

**Note:**  $p(\cdot)$  corresponds to a shadow price and (P3) to Kelly's willingness to pay type control.

# Further study

- How to generate signals using only data obtained by the network
- Analysis of a model that includes delay of signal